

A Platform for Comparing CDS Discovering Algorithms

M.S. PROJECT REPORT
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Outline

- **Problem and Goal**
 - Develop a platform for CDS algorithms comparison
- **Background**
 - DS, CDS, MCDS
- **Three Algorithms**
- **Design Guidelines**
 - OS-independent, MVC, OO design
- **Design Architecture**
 - Packages, UML
- **Demo**
- **Conclusion**

Problem

- **Connected Dominating Sets (CDS)**
 - Important in many applications, e.g., wireless sensor network
- **Too many existing algorithms**
 - Famous ones: Alzoubi's, Wu Jie's, R-value (Dr. Li's group)
 - Hard to compare their performances as they are implemented in different platforms and languages
- **Question:**
 - How to provide a research facility to easily simulate, test, and compare various CDS algorithms?

Goal

Background

- **Dominating sets:**
 - A **dominating set** for a graph $G = (V, E)$ is a subset D of V such that every vertex not in D is joined to at least one member of D by some edge.

(a)

(b)

(c)

- **Connected Dominating sets**
 - A connected dominating set (CDS) is a subset of the nodes such that it forms a DS and all the nodes in the DS are connected.

Background

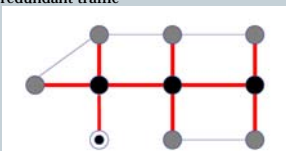
- **Minimum Connected Dominating sets**
 - A **minimum connected dominating set (MCDS)** of a graph G is a CDS with the smallest number of nodes.

CDS

MCDS

Background

- One example of CDS application
 - CDS is used as a virtual backbone in wireless networks.
 - Network broadcasting:
 - Only nodes in CDS relay messages
 - Reduce communication cost
 - Reduce redundant traffic



Three MCDS Algorithms

- Alzoubi et al. Algorithm
- Wu's Algorithm
- Our R-value Algorithm

Alzoubi's Algorithm

1. Construct a spanning tree

- Assign level 0 to root node v
- Propagate message and assign level for each node

2. Find MIS nodes in even levels

- Mark nodes to **black** based on its rank, others are **grey**
- Rank priority: $\langle \text{level}, \text{id} \rangle$

3. Find connectors in odd levels to form a CDS

- A connector should connect to 2 unconnected black nodes, mark as **blue**
- All **black** and **blue** nodes are **dominators**

Wujie's Algorithm

Open neighbor set $N(v)$, close neighbor set $N[v]$

1. Marking process

- All nodes are initially **white**
- Check each node v : if two neighbors of v are unconnected, mark v **black**
- All **black** nodes form a CDS

2. Pruning process

- If u, v are **black**, and $N[v] \subset N[u]$, ~~unmark~~ v
- If u, v, w are **black**, and $N[v] \subset N[u] \cup N[w]$, ~~unmark~~ v

r-CDS algorithm

Definition of r-value and deg

- $r(u)$ = number of r -hop neighbors - $d(u)$
- $\text{deg}(u)$ = number of **white** neighbors
- Rank priority is triplet $\langle r, \text{deg}, \text{id} \rangle$

1. Marking step:

- Check each node v , mark it **black** if it has lowest rank among neighbors
- Mark black node's neighbors **grey**

2. Connecting step - find connectors to form CDS

- Mark a **grey** node **black** if it connects 2 unconnected black nodes
- Mark two **grey** nodes **black** if they together connects 2 unconnected black nodes

r-CDS algorithm

3. Pruning step

- For two **black** nodes u, v , mark u **grey** if
 - $N[u] \subset N[v]$ or
 - $N[u] = N[v]$ and $\text{id}(u) < \text{id}(v)$
- Mark u **grey** if u has 2 **black** neighbors v, w , such that:
 - $N(u) \subset N(v) \cup N(w) - \{u\}$ or
 - $N(u) - \{v, w\} = N(v) \cup N(w) - \{u, v, w\}$, and $\text{id}(u)$ is minimum
- Mark u **grey** if u has 3 **black** neighbors x, y, z , such that:
 - $N(u) \subset N(x) \cup N(y) \cup N(z) - \{u\}$ or
 - $N(u) - \{x, y, z\} = N(x) \cup N(y) \cup N(z) - \{x, y, z\}$, and $\text{id}(u)$ is minimum

